

# Randomized SVD, CUR Decomposition, and SPSD Matrix Approximation

**Shusen Wang**

# Outline

- CX Decomposition & Approximate SVD
- CUR Decomposition
- SPSD Matrix Approximation

# CX Decomposition

- Given any matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$
- The CX decomposition of  $\mathbf{A}$ 
  1. Sketching:  $\mathbf{C} = \mathbf{A}\mathbf{P} \in \mathbb{R}^{m \times c}$
  2. Find  $\mathbf{X}$  such that  $\mathbf{A} \approx \mathbf{C}\mathbf{X}$ 
    - E.g.  $\mathbf{X}^* = \operatorname{argmin}_{\mathbf{X}} \|\mathbf{A} - \mathbf{C}\mathbf{X}\|_{\mathbf{F}}^2 = \mathbf{C}^\dagger \mathbf{A}$
    - It costs  $O(mnc)$

# CX Decomposition

- Let the sketching matrix  $\mathbf{P} \in \mathbb{R}^{n \times c}$  be defined in the table.
- $$\min_{\text{rank}(\mathbf{X}) \leq k} \|\mathbf{A} - \mathbf{CX}\|_F^2 \leq (1 + \epsilon) \|\mathbf{A} - \mathbf{A}_k\|_F^2$$

|          | Uniform sampling   | Leverage score sampling                                    | Gaussian projection                | SRHT  | Count sketch                             |
|----------|--|--|------------------------------------|---|--|
| $c \geq$ | $O\left(\nu k \left(\log k + \frac{1}{\epsilon}\right)\right)$ | $O\left(k \left(\log k + \frac{1}{\epsilon}\right)\right)$ | $O\left(\frac{k}{\epsilon}\right)$ | $O\left((k + \log n) \left(\log k + \frac{1}{\epsilon}\right)\right)$ | $O\left(k^2 + \frac{k}{\epsilon}\right)$ |

$\nu$  is the column coherence of  $\mathbf{A}_k$

# CX Decomposition $\Leftrightarrow$ Approximate SVD


- CX decomposition  $\Leftrightarrow$  approximate SVD

$$A \approx CX$$

# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD

$$\mathbf{A} \approx \mathbf{CX} = \mathbf{U}_C \mathbf{\Sigma}_C \mathbf{V}_C^T \mathbf{X}$$



SVD:  $\mathbf{C} = \mathbf{U}_C \mathbf{\Sigma}_C \mathbf{V}_C^T \in \mathbb{R}^{m \times c}$

Time cost:  $O(mc^2)$

# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD

$$\mathbf{A} \approx \mathbf{C}\mathbf{X} = \mathbf{U}_C \boldsymbol{\Sigma}_C \mathbf{V}_C^T \mathbf{X} = \mathbf{U}_C \mathbf{Z}$$

SVD:  $\mathbf{C} = \mathbf{U}_C \boldsymbol{\Sigma}_C \mathbf{V}_C^T \in \mathbb{R}^{m \times c}$

Let  $\boldsymbol{\Sigma}_C \mathbf{V}_C^T \mathbf{X} = \mathbf{Z} \in \mathbb{R}^{c \times n}$

Time cost:  $O(mc^2 + nc^2)$

# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD

$$\mathbf{A} \approx \mathbf{C}\mathbf{X} = \mathbf{U}_C \boldsymbol{\Sigma}_C \mathbf{V}_C^T \mathbf{X} = \mathbf{U}_C \mathbf{Z} = \mathbf{U}_C \mathbf{U}_Z \boldsymbol{\Sigma}_Z \mathbf{V}_Z^T$$

Let  $\boldsymbol{\Sigma}_C \mathbf{V}_C^T \mathbf{X} = \mathbf{Z} \in \mathbb{R}^{c \times n}$

SVD:  $\mathbf{C} = \mathbf{U}_C \boldsymbol{\Sigma}_C \mathbf{V}_C^T \in \mathbb{R}^{m \times c}$

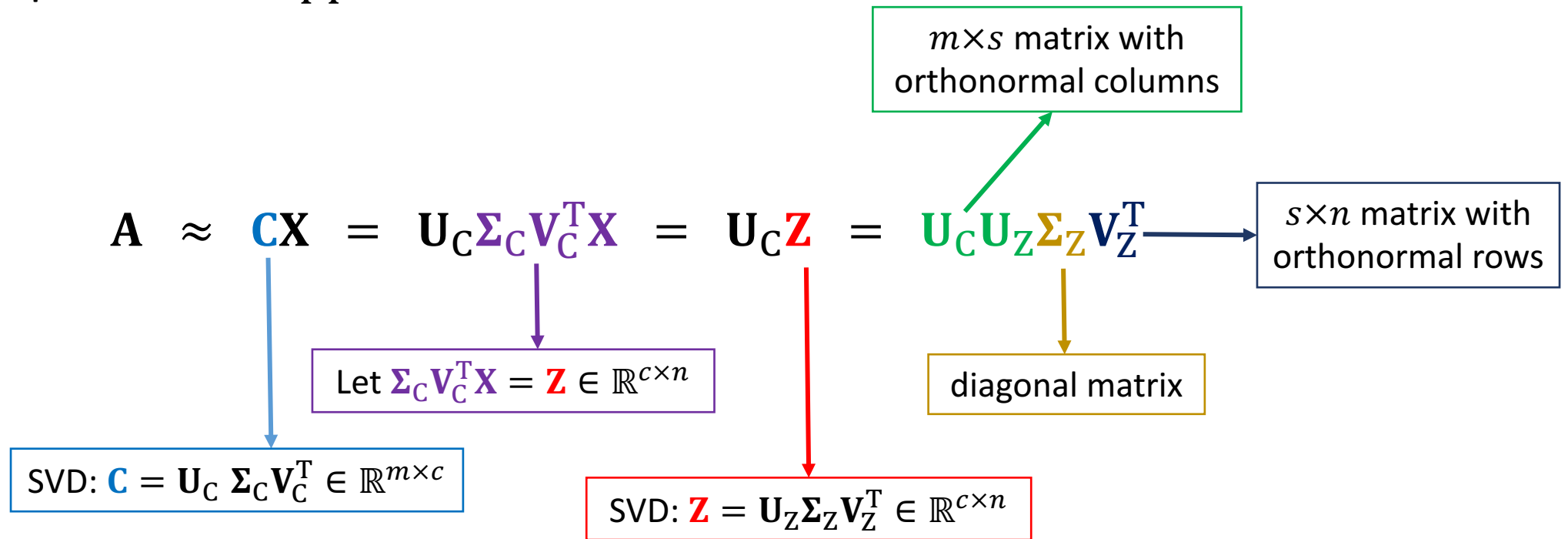
SVD:  $\mathbf{Z} = \mathbf{U}_Z \boldsymbol{\Sigma}_Z \mathbf{V}_Z^T \in \mathbb{R}^{c \times n}$

Time cost:  $O(mc^2 + nc^2 + nc^2)$



# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD



Time cost:  $O(mc^2 + nc^2 + nc^2 + mc^2)$

# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD
- Done! Approximate rank  $c$  SVD:  $\mathbf{A} \approx (\mathbf{U}_C \mathbf{U}_Z) \mathbf{\Sigma}_Z \mathbf{V}_Z^T$

$$\mathbf{A} \approx \mathbf{C}\mathbf{X} = \mathbf{U}_C \mathbf{\Sigma}_C \mathbf{V}_C^T \mathbf{X} = \mathbf{U}_C \mathbf{Z} = \mathbf{U}_C \mathbf{U}_Z \mathbf{\Sigma}_Z \mathbf{V}_Z^T$$

The diagram illustrates the decomposition of the product matrix  $\mathbf{U}_C \mathbf{U}_Z$  into three components:

- A green box labeled " $m \times s$  matrix with orthonormal columns" has a green arrow pointing to the  $\mathbf{U}_C \mathbf{U}_Z$  term in the equation.
- A blue box labeled " $s \times n$  matrix with orthonormal rows" has a blue arrow pointing to the  $\mathbf{V}_Z^T$  term in the equation.
- A yellow box labeled "diagonal matrix" has a yellow arrow pointing to the  $\mathbf{\Sigma}_Z$  term in the equation.

Time cost:  $O(mc^2 + nc^2 + nc^2 + mc^2) = O(mc^2 + nc^2)$

# CX Decomposition $\Leftrightarrow$ Approximate SVD

- CX decomposition  $\Leftrightarrow$  approximate SVD
- Given  $\mathbf{A} \in \mathbb{R}^{m \times n}$  and  $\mathbf{C} \in \mathbb{R}^{m \times c}$ , the approximate SVD costs
  - $O(mnc)$  time
  - $O(mc + nc)$  memory

# CX Decomposition

- The CX decomposition of  $\mathbf{A} \in \mathbb{R}^{m \times n}$ 
  - Optimal solution:  $\mathbf{X}^* = \operatorname{argmin}_{\mathbf{X}} \|\mathbf{A} - \mathbf{CX}\|_{\text{F}}^2 = \mathbf{C}^\dagger \mathbf{A}$
  - How to make it more efficient?

# CX Decomposition

- The CX decomposition of  $\mathbf{A} \in \mathbb{R}^{m \times n}$

- Optimal solution:  $\mathbf{X}^* = \operatorname{argmin}_{\mathbf{X}} \|\mathbf{A} - \mathbf{CX}\|_{\text{F}}^2 = \mathbf{C}^\dagger \mathbf{A}$

- How to make it more efficient?

A regression problem!

# Fast CX Decomposition

- Fast CX [Drineas, Mahoney, Muthukrishnan, 2008][Clarkson & Woodruff, 2013]
  - Draw another sketching matrix  $\mathbf{S} \in \mathbb{R}^{m \times s}$
  - Compute  $\tilde{\mathbf{X}} = \operatorname{argmin}_{\mathbf{X}} \|\mathbf{S}^T (\mathbf{A} - \mathbf{C}\mathbf{X})\|_{\text{F}}^2 = (\mathbf{S}^T \mathbf{C})^\dagger (\mathbf{S}^T \mathbf{A})$
  - Time cost:  $O(ncs) + \text{TimeOfSketch}$
  - When  $s = \tilde{O}(c/\epsilon)$ ,
$$\|\mathbf{A} - \mathbf{C}\tilde{\mathbf{X}}\|_{\text{F}}^2 \leq (1 + \epsilon) \cdot \min_{\mathbf{X}} \|\mathbf{A} - \mathbf{C}\mathbf{X}\|_{\text{F}}^2$$

# Outline

- CX Decomposition & Approximate SVD
- **CUR Decomposition**
- SPSD Matrix Approximation

# CUR Decomposition

- Sketching
  - $\mathbf{C} = \mathbf{A}\mathbf{P}_C \in \mathbb{R}^{m \times c}$
  - $\mathbf{R} = \mathbf{P}_R^T \mathbf{A} \in \mathbb{R}^{r \times n}$
- Find  $\mathbf{U}$  such that  $\mathbf{CUR} \approx \mathbf{A}$
- $\mathbf{CUR} \Leftrightarrow$  Approximate SVD
  - In the same way as “ $\mathbf{CX} \Leftrightarrow$  Approximate SVD”



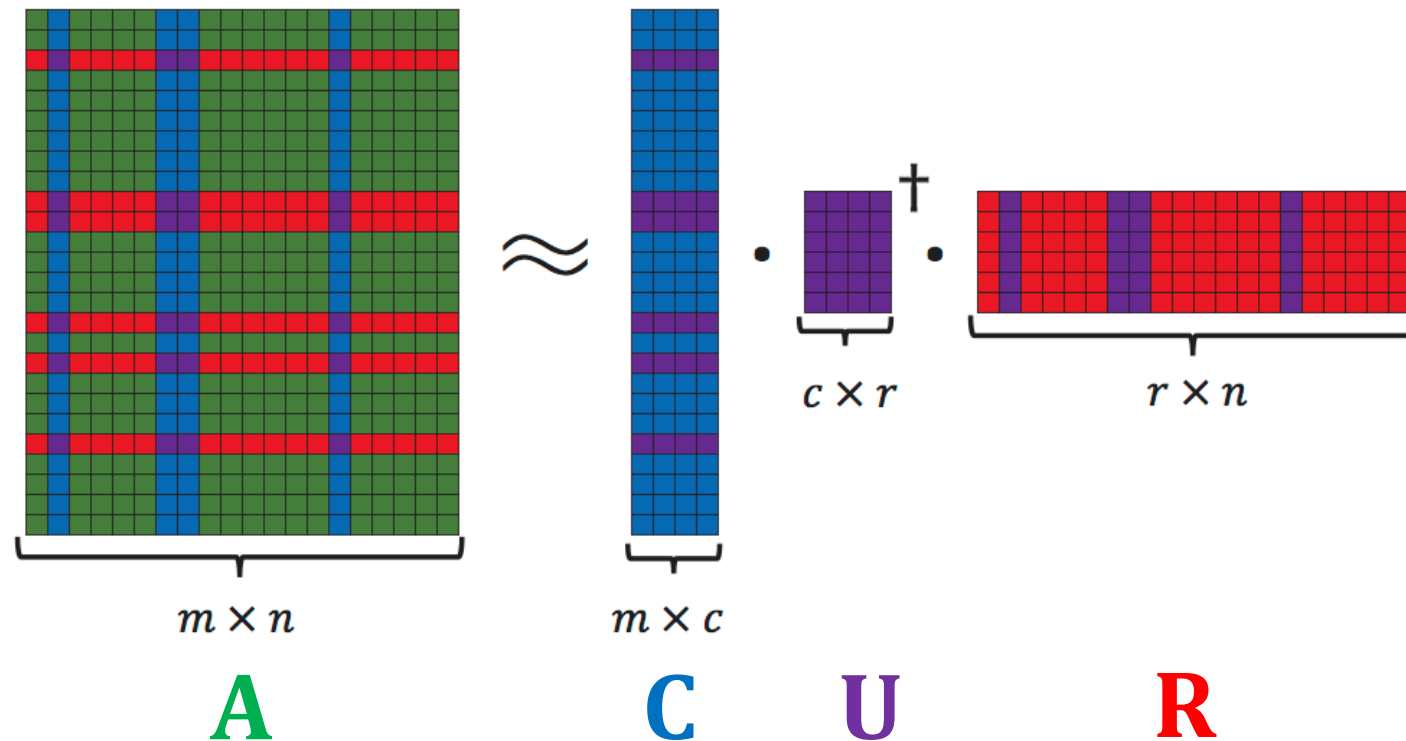
# CUR Decomposition

- Sketching
  - $\mathbf{C} = \mathbf{A}\mathbf{P}_C \in \mathbb{R}^{m \times c}$
  - $\mathbf{R} = \mathbf{P}_R^T \mathbf{A} \in \mathbb{R}^{r \times n}$
- Find  $\mathbf{U}$  such that  $\mathbf{CUR} \approx \mathbf{A}$
- $\mathbf{CUR} \Leftrightarrow$  Approximate SVD
  - In the same way as “ $\mathbf{CX} \Leftrightarrow$  Approximate SVD”
- 3 types of  $\mathbf{U}$

# CUR Decomposition

- Type 1 [Drineas, Mahoney, Muthukrishnan, 2008]:

$$U = (P_R^T A P_C)^{\dagger}$$



# CUR Decomposition

- Type 1 [Drineas, Mahoney, Muthukrishnan, 2008]:

$$\mathbf{U} = (\mathbf{P}_R^T \mathbf{A} \mathbf{P}_C)^{\dagger}$$

- Recall the fast CX decomposition

$$\mathbf{A} \approx \mathbf{C} \tilde{\mathbf{X}} = \mathbf{C} (\mathbf{P}_R^T \mathbf{C})^{\dagger} (\mathbf{P}_R^T \mathbf{A})$$

# CUR Decomposition

- Type 1 [Drineas, Mahoney, Muthukrishnan, 2008]:

$$\mathbf{U} = (\mathbf{P}_R^T \mathbf{A} \mathbf{P}_C)^{\dagger}$$

- Recall the fast CX decomposition

$$\mathbf{A} \approx \mathbf{C}\tilde{\mathbf{X}} = \mathbf{C}(\mathbf{P}_R^T \mathbf{C})^{\dagger} (\mathbf{P}_R^T \mathbf{A}) = \mathbf{C}\mathbf{U}\mathbf{R}$$

# CUR Decomposition

- Type 1 [Drineas, Mahoney, Muthukrishnan, 2008]:

$$\mathbf{U} = (\mathbf{P}_R^T \mathbf{A} \mathbf{P}_C)^{\dagger}$$

- Recall the fast CX decomposition

$$\mathbf{A} \approx \mathbf{C} \tilde{\mathbf{X}} = \mathbf{C} (\mathbf{P}_R^T \mathbf{C})^{\dagger} (\mathbf{P}_R^T \mathbf{A}) = \mathbf{C} \mathbf{U} \mathbf{R}$$

- They're equivalent:  $\mathbf{C} \tilde{\mathbf{X}} = \mathbf{C} \mathbf{U} \mathbf{R}$

# CUR Decomposition

- **Type 1** [Drineas, Mahoney, Muthukrishnan, 2008]:

$$\mathbf{U} = (\mathbf{P}_R^T \mathbf{A} \mathbf{P}_C)^{\dagger}$$

- Recall the fast CX decomposition

$$\mathbf{A} \approx \mathbf{C} \tilde{\mathbf{X}} = \mathbf{C} (\mathbf{P}_R^T \mathbf{C})^{\dagger} (\mathbf{P}_R^T \mathbf{A}) = \mathbf{C} \mathbf{U} \mathbf{R}$$

- They're equivalent:  $\mathbf{C} \tilde{\mathbf{X}} = \mathbf{C} \mathbf{U} \mathbf{R}$

- Require  $c = \tilde{O}\left(\frac{k}{\epsilon}\right)$  and  $r = \tilde{O}\left(\frac{c}{\epsilon}\right)$  such that

$$\|\mathbf{A} - \mathbf{C} \mathbf{U} \mathbf{R}\|_F^2 \leq (1 + \epsilon) \|\mathbf{A} - \mathbf{A}_k\|_F^2$$

# CUR Decomposition

- Type 1 [Drineas, Mahoney, Muthukrishnan, 2008]:

$$\mathbf{U} = (\mathbf{P}_R^T \mathbf{A} \mathbf{P}_C)^{\dagger}$$

- Efficient
  - $O(rc^2) + \text{TimeOfSketch}$
- Loose bound
  - Sketch size  $\propto \epsilon^{-2}$
- Bad empirical performance

# CUR Decomposition

- Type 2: Optimal CUR

$$\mathbf{U}^* = \min_{\mathbf{U}} \|\mathbf{A} - \mathbf{CUR}\|_F^2 = \mathbf{C}^\dagger \mathbf{A} \mathbf{R}^\dagger$$



# CUR Decomposition

- Type 2: Optimal CUR

$$\mathbf{U}^* = \min_{\mathbf{U}} \|\mathbf{A} - \mathbf{CUR}\|_F^2 = \mathbf{C}^\dagger \mathbf{A} \mathbf{R}^\dagger$$

- Theory [W & Zhang, 2013], [Boutsidis & Woodruff, 2014]:

- $\mathbf{C}$  and  $\mathbf{R}$  are selected by the adaptive sampling algorithm

- $c = O\left(\frac{k}{\epsilon}\right)$  and  $r = O\left(\frac{k}{\epsilon}\right)$

- $\|\mathbf{A} - \mathbf{CUR}\|_F^2 \leq (1 + \epsilon) \|\mathbf{A} - \mathbf{A}_k\|_F^2$

# CUR Decomposition

- Type 2: Optimal CUR

$$\mathbf{U}^* = \min_{\mathbf{U}} \|\mathbf{A} - \mathbf{CUR}\|_F^2 = \mathbf{C}^\dagger \mathbf{A} \mathbf{R}^\dagger$$

- Inefficient
  - $O(mnc)$  + TimeOfSketch

# CUR Decomposition

- Type 3: Fast CUR [W, Zhang, Zhang, 2015]
  - Draw 2 sketching matrices  $\mathbf{S}_C$  and  $\mathbf{S}_R$
  - Solve the problem

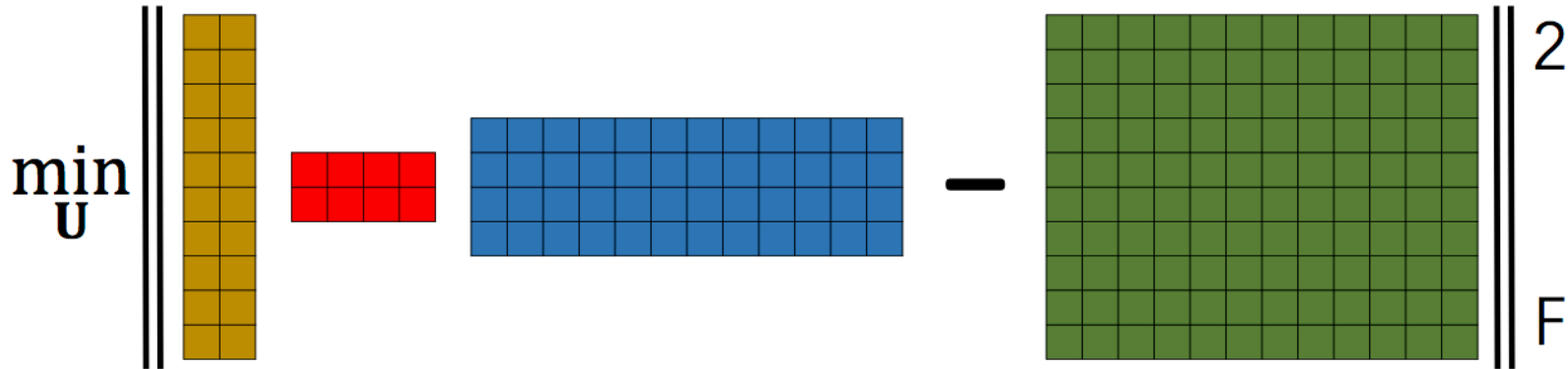
$$\tilde{\mathbf{U}} = \min_{\mathbf{U}} \left\| \mathbf{S}_C^T (\mathbf{A} - \mathbf{CUR}) \mathbf{S}_R \right\|_F^2 = (\mathbf{S}_C^T \mathbf{C})^\dagger (\mathbf{S}_C^T \mathbf{A} \mathbf{S}_R) (\mathbf{R} \mathbf{S}_R)^\dagger$$

- Intuition?

# CUR Decomposition

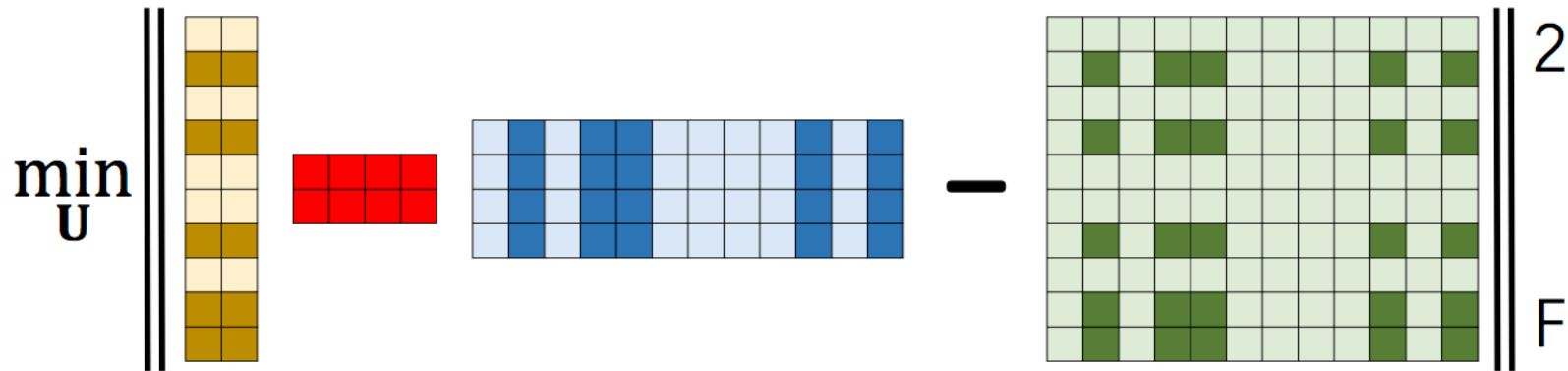
- The optimal  $\mathbf{U}$  matrix is obtained by the optimization problem

$$\mathbf{U}^* = \min_{\mathbf{U}} \|\mathbf{CUR} - \mathbf{A}\|_F^2$$



# CUR Decomposition

- Approximately solve the optimization problem, e.g. by column selection



# CUR Decomposition

- Solve the small scale problem

$$\min_U \left\| \begin{array}{c} \text{yellow grid} \\ \text{red grid} \end{array} \begin{array}{c} \text{blue grid} \end{array} - \begin{array}{c} \text{green grid} \end{array} \right\|_F^2$$

The diagram illustrates the CUR decomposition of a matrix. It shows a minimization problem over matrix U. The expression is:  $\min_U \left\| \begin{array}{c} \text{yellow grid} \\ \text{red grid} \end{array} \begin{array}{c} \text{blue grid} \end{array} - \begin{array}{c} \text{green grid} \end{array} \right\|_F^2$ . The yellow grid is 5x2, the red grid is 2x3, the blue grid is 5x3, and the green grid is 5x3. The norm is the Frobenius norm, denoted by  $\| \cdot \|_F$ .

# CUR Decomposition

- Type 3: Fast CUR [W, Zhang, Zhang, 2015]

- Draw 2 sketching matrices  $\mathbf{S}_C \in \mathbb{R}^{m \times s_c}$  and  $\mathbf{S}_R \in \mathbb{R}^{n \times s_r}$
- Solve the problem

$$\tilde{\mathbf{U}} = \min_{\mathbf{U}} \left\| \mathbf{S}_C^T (\mathbf{A} - \mathbf{CUR}) \mathbf{S}_R \right\|_F^2 = (\mathbf{S}_C^T \mathbf{C})^\dagger (\mathbf{S}_C^T \mathbf{A} \mathbf{S}_R) (\mathbf{R} \mathbf{S}_R)^\dagger$$

- Theory

- $s_c = O\left(\frac{c}{\epsilon}\right)$  and  $s_r = O\left(\frac{r}{\epsilon}\right)$
- $\left\| \mathbf{A} - \mathbf{C}\tilde{\mathbf{U}}\mathbf{R} \right\|_F^2 \leq (1 + \epsilon) \cdot \min_{\mathbf{U}} \left\| \mathbf{A} - \mathbf{CUR} \right\|_F^2$

# CUR Decomposition

- Type 3: Fast CUR [W, Zhang, Zhang, 2015]

- Draw 2 sketching matrices  $\mathbf{S}_C \in \mathbb{R}^{m \times s_c}$  and  $\mathbf{S}_R \in \mathbb{R}^{n \times s_r}$
- Solve the problem

$$\tilde{\mathbf{U}} = \min_{\mathbf{U}} \left\| \mathbf{S}_C^T (\mathbf{A} - \mathbf{CUR}) \mathbf{S}_R \right\|_F^2 = (\mathbf{S}_C^T \mathbf{C})^\dagger (\mathbf{S}_C^T \mathbf{A} \mathbf{S}_R) (\mathbf{R} \mathbf{S}_R)^\dagger$$

- Efficient
  - $O(s_c s_r (c + r)) + \text{TimeOfSketch}$
- Good empirical performance



**A:**  
 $m = 1920$   
 $n = 1168$

**C and R:**

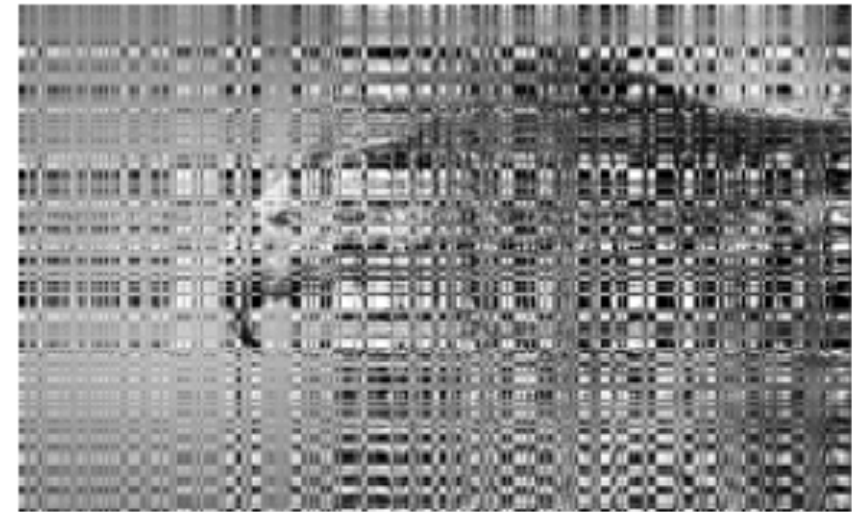
- $c = r = 100$
- uniform sampling



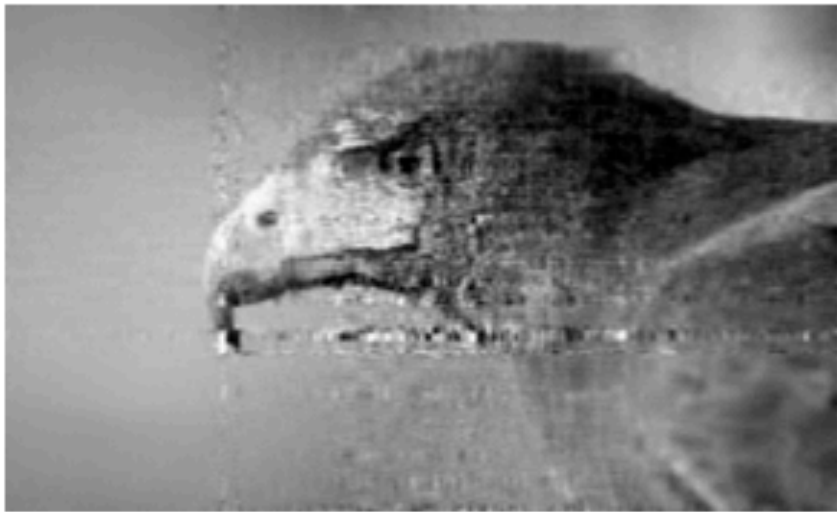
Original



Type 2: Optimal CUR



Type 1: Fast CX



Type 3: Fast CUR  
 $s_c = 2c, \quad s_r = 2r$



Type 3: Fast CUR  
 $s_c = 4c, \quad s_r = 4r$

# Conclusions

- Approximate truncated SVD
  - CX decomposition
  - CUR decomposition (3 types)
- Fast CUR is the best

# Outline

- CX Decomposition & Approximate SVD
- CUR Decomposition
- **SPSD Matrix Approximation**

# Motivation 1: Kernel Matrix

- Given  $n$  samples  $\mathbf{x}_1, \dots, \mathbf{x}_n \in \mathbb{R}^d$  and kernel function  $\kappa(\cdot, \cdot)$ .
- E.g. Gaussian RBF kernel

$$\kappa(\mathbf{x}_i, \mathbf{x}_j) = \exp\left(-\frac{\|\mathbf{x}_i - \mathbf{x}_j\|_2^2}{\sigma^2}\right).$$

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- E.g. Gaussian RBF kernel

$$\kappa(\mathbf{x}_i, \mathbf{x}_j) = \exp\left(-\frac{\|\mathbf{x}_i - \mathbf{x}_j\|_2^2}{\sigma^2}\right).$$

- Computing the kernel matrix  $\mathbf{K} \in \mathbb{R}^{n \times n}$ 
  - where  $k_{ij} = \kappa(\mathbf{x}_i, \mathbf{x}_j)$
  - costs  $O(n^2d)$  time

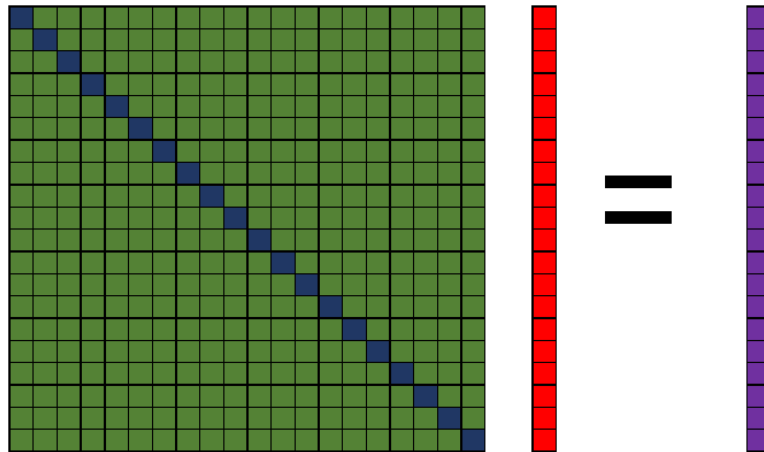
# Motivation 2: Matrix Inversion

- Solve the linear system

$$(\mathbf{K} + \alpha \mathbf{I}_n) \mathbf{w} = \mathbf{y}$$

to find  $\mathbf{w} \in \mathbb{R}^n$ .

- $\mathbf{K} \in \mathbb{R}^{n \times n}$  is the kernel matrix
- $\mathbf{y} = [y_1, \dots, y_n] \in \mathbb{R}^n$  contains the labels



# Motivation 2: Matrix Inversion

- Solve the linear system

$$(\mathbf{K} + \alpha \mathbf{I}_n) \mathbf{w} = \mathbf{y}$$

to find  $\mathbf{w} \in \mathbb{R}^n$ .

- Solution:  $\mathbf{w}^* = (\mathbf{K} + \alpha \mathbf{I}_n)^{-1} \mathbf{y}$

# Motivation 2: Matrix Inversion

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- Solution:  $\mathbf{w}^* = (\mathbf{K} + \alpha \mathbf{I}_n)^{-1} \mathbf{y}$
- It costs
  - $O(n^3)$  time
  - $O(n^2)$  memory.



# Motivation 2: Matrix Inversion

- Solve the linear system

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- Solution:  $\mathbf{w}^* = (\mathbf{K} + \alpha \mathbf{I}_n)^{-1} \mathbf{y}$
- It costs
  - $O(n^3)$  time
  - $O(n^2)$  memory.
- Performed by
  - Kernel ridge regression
  - Least squares kernel SVM

# Motivation 3: Eigenvalue Decomposition

- Find the top  $k$  ( $\ll n$ ) eigenvectors of  $\mathbf{K}$ .
- It costs
  - $\tilde{O}(n^2 k)$  time
  - $O(n^2)$  memory.

# Motivation 3: Eigenvalue Decomposition

- Find the top  $k$  ( $\ll n$ ) eigenvectors of  $\mathbf{K}$ .
- It costs
  - $\tilde{O}(n^2k)$  time
  - $O(n^2)$  memory.
- Performed by
  - Kernel PCA ( $k$  is the target rank)
  - Manifold learning ( $k$  is the target rank)

# Computational Challenges

- Time costs
  - Computing kernel matrix:  $O(n^2d)$
  - Matrix inversion:  $O(n^3)$
  - Rank  $k$  eigenvalue decomposition:  $O(n^2k)$

# Computational Challenges

- Time costs

- Computing kernel matrix:  $O(n^2 d)$
- Matrix inversion:  $O(n^3)$
- Rank  $k$  eigenvalue decomposition:  $O(n^2 k)$

**At least quadratic time!**

# Computational Challenges

- Time costs
  - Computing kernel matrix:  $O(n^2d)$
  - Matrix inversion:  $O(n^3)$
  - Rank  $k$  eigenvalue decomposition:  $O(n^2k)$
- Memory costs
  - Inversion and eigenvalue decomposition:  $O(n^2)$

# Computational Challenges

- Time costs
  - Computing kernel matrix:  $O(n^2d)$
  - Matrix inversion:  $O(n^3)$
  - Rank  $k$  eigenvalue decomposition:  $O(n^2k)$
- Memory costs
  - Inversion and eigenvalue decomposition:  $O(n^2)$
  - Because
    - the numerical algorithms are pass-inefficient
    - → form  $\mathbf{K}$  and keep it in memory

# Computational Challenges

- Time costs
  - Computing kernel matrix:  $O(n^2d)$
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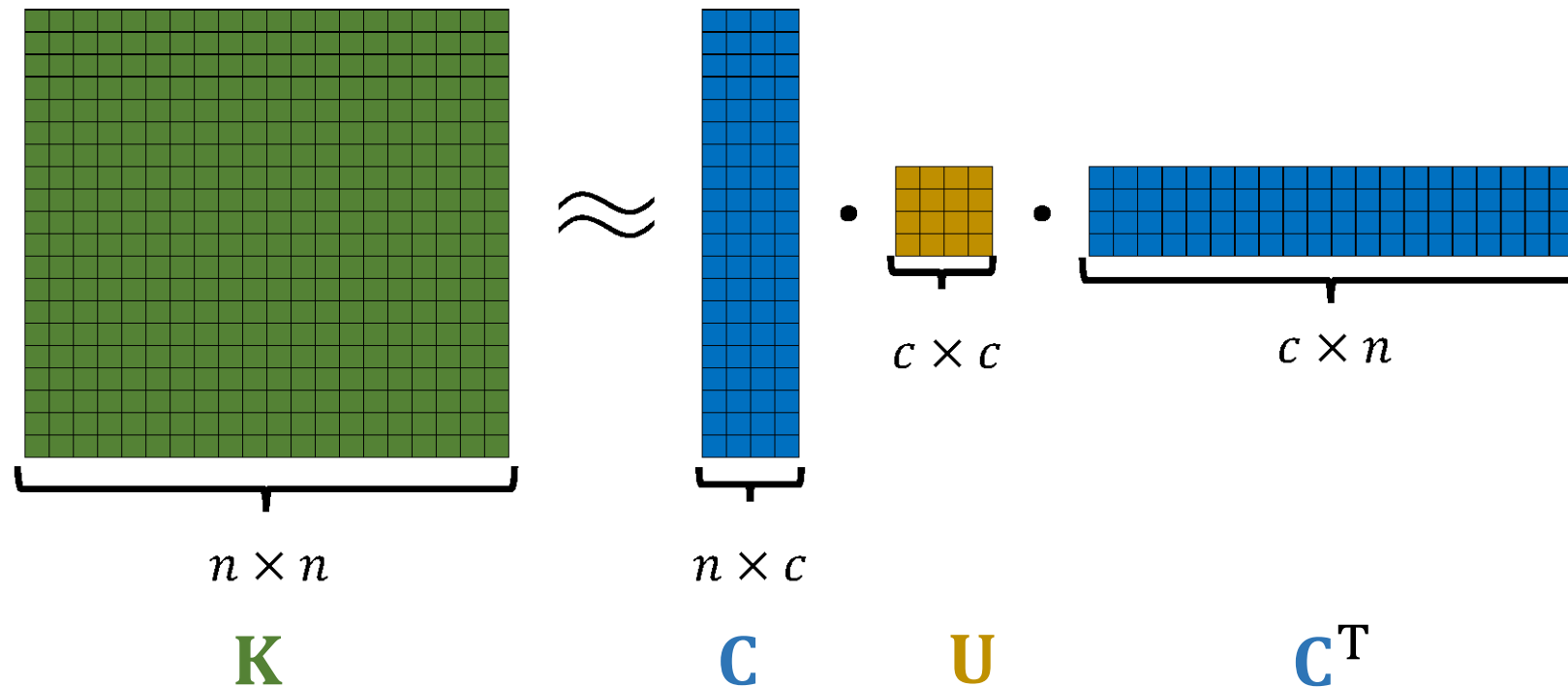
When  $n = 10^5$ , the  $n \times n$  matrix costs 80GB memory!



# How to Speedup?

- Efficiently form the low-rank approximation

$$\mathbf{K} \approx \mathbf{C} \mathbf{U} \mathbf{C}^T$$



# How to Speedup?

- Efficiently form the low-rank approximation

$$\mathbf{K} \approx \mathbf{C} \mathbf{U} \mathbf{C}^T$$

- Equivalent  $\mathbf{K} \approx \mathbf{L} \mathbf{L}^T$

# Efficient Matrix Inversion

- Solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ :

$$\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$$

# Efficient Matrix Inversion

- Approximately solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ :
  - Replace  $\mathbf{K}$  by  $\mathbf{LL}^T$ :  $\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y} \approx (\mathbf{LL}^T + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$

# Efficient Matrix Inversion

- Approximately solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ 
  - Replace  $\mathbf{K}$  by  $\mathbf{L}\mathbf{L}^T$ :  $\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y} \approx (\mathbf{L}\mathbf{L}^T + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$
  - Expand the **inversion** by the Woodbury identity

# Efficient Matrix Inversion

- Approximately solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ 
  - Replace  $\mathbf{K}$  by  $\mathbf{L}\mathbf{L}^T$ :  $\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y} \approx (\mathbf{L}\mathbf{L}^T + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$
  - Expand the inversion by the Woodbury identity

$$(\mathbf{A} + \mathbf{BCD})^{-1} = \mathbf{A}^{-1} - \mathbf{A}^{-1}\mathbf{B}(\mathbf{C}^{-1} + \mathbf{DA}^{-1}\mathbf{B})^{-1}\mathbf{DA}^{-1}$$

# Efficient Matrix Inversion

- Approximately solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ 
  - Replace  $\mathbf{K}$  by  $\mathbf{L}\mathbf{L}^T$ :  $\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y} \approx (\mathbf{L}\mathbf{L}^T + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$
  - Expand the inversion by the Woodbury identity
$$\mathbf{w}^* \approx \alpha^{-1}\mathbf{y} + \alpha^{-1}\mathbf{L}(\alpha\mathbf{I} + \mathbf{L}^T\mathbf{L})^{-1}\mathbf{L}^T\mathbf{y}$$

# Efficient Matrix Inversion

- Approximately solve the linear system  $(\mathbf{K} + \alpha\mathbf{I}_n)\mathbf{w} = \mathbf{y}$ 
  - Replace  $\mathbf{K}$  by  $\mathbf{L}\mathbf{L}^T$ :  $\mathbf{w}^* = (\mathbf{K} + \alpha\mathbf{I}_n)^{-1}\mathbf{y} \approx (\mathbf{L}\mathbf{L}^T + \alpha\mathbf{I}_n)^{-1}\mathbf{y}$
  - Expand the inversion by the Woodbury identity

$$\mathbf{w}^* \approx \alpha^{-1}\mathbf{y} + \alpha^{-1}\mathbf{L}(\alpha\mathbf{I} + \mathbf{L}^T\mathbf{L})^{-1}\mathbf{L}^T\mathbf{y}$$

- Time cost:  $O(nc^2)$

Linear in  $n$ , much better than  $O(n^3)$



# Efficient Eigenvalue Decomposition

- Approximately compute the  $k$ -eigenvalue decomposition of  $\mathbf{K}$ 
  - SVD:  $\mathbf{L} = \mathbf{U}_L \mathbf{\Sigma}_L \mathbf{V}_L$
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  - Approximate  $k$ - eigenvalue decomposition of  $\mathbf{K}$ 
    - eigenvectors: the first  $k$  vectors in  $\mathbf{U}_L$
- Time cost:  $O(nc^2)$

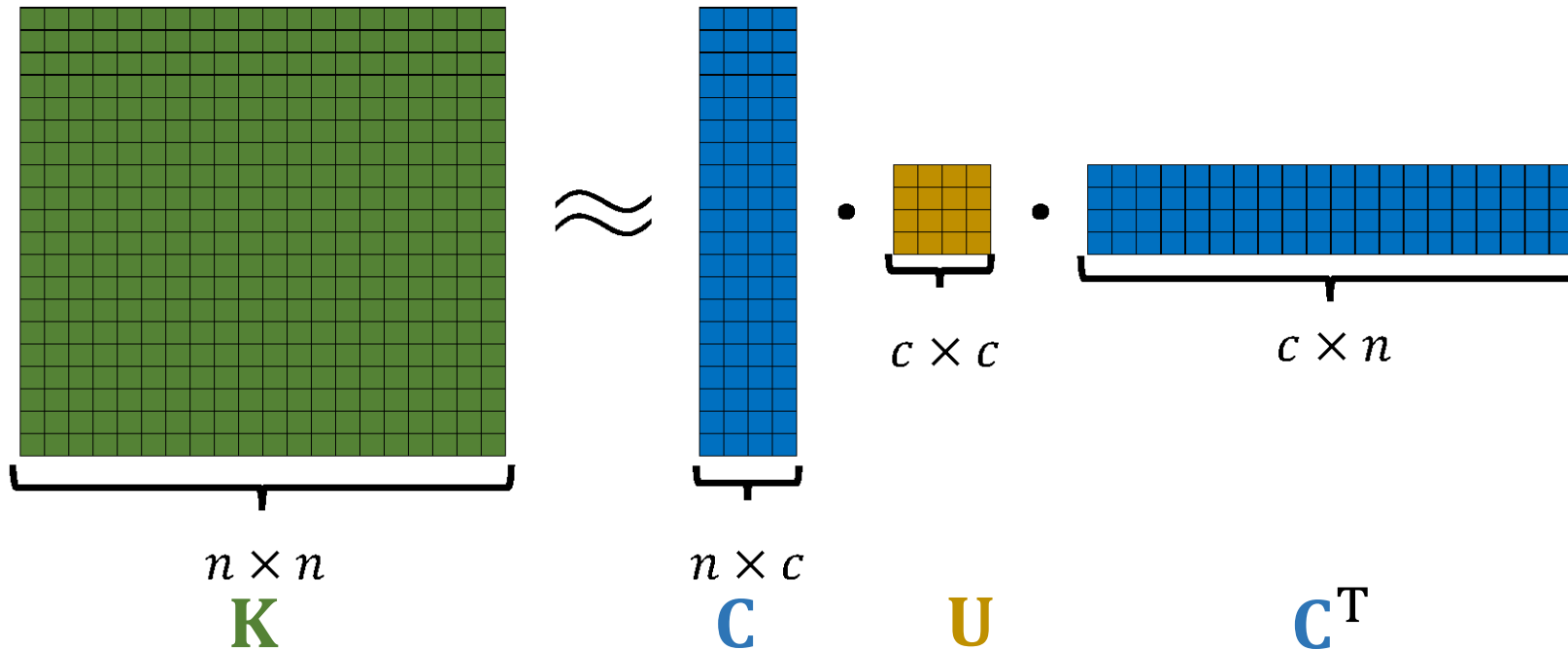
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    - eigenvectors: the first  $k$  vectors in  $\mathbf{U}_L$
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  - Much lower than  $\tilde{O}(n^2k)$

# Sketching Based Models

- How to find such an approximation?

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  - $\mathbf{S} \in \mathbb{R}^{n \times c}$  can be column selection or random projection matrix
- Three methods:
  - The prototype model [HMT11, WZ13, WLZ16]
  - The fast model [WZZ15]
  - The Nyström method [WS15, GM13]

# The Prototype Model

- Objective:  $\mathbf{K} \approx \mathbf{CUC}^T$

- Minimize the approximation error by

$$\mathbf{U}^* = \operatorname{argmin}_{\mathbf{U}} \left\| \mathbf{K} - \mathbf{CUC}^T \right\|_F^2 = \mathbf{C}^+ \mathbf{K} (\mathbf{C}^+)^T.$$

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Extension of the random SVD to SPSD matrix [HMT11]



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- Time:  $O(n^2c)$

- The time complexity is nearly the same to the  $k$ -eigenvalue decomposition.
- It is much faster than the  $k$ - eigenvalue decomposition in practice.

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- #Passes: one

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- Time:  $O(n^2c)$
- #Passes: one
- Memory:  $O(nc)$ 
  - Put  $k_{ij}$  in memory only when it is visited
  - Keep  $\mathbf{C}^\dagger$  in memory

# The Prototype Model

- Error Bound
  - $k \ll n$  is arbitrary integer
  - $\mathbf{P}$  samples  $c = O\left(\frac{k}{\epsilon}\right)$  columns by adaptive sampling
  - $\mathbb{E} \left\| \mathbf{K} - \mathbf{C}\mathbf{U}^*\mathbf{C}^T \right\|_F^2 \leq (1 + \epsilon) \left\| \mathbf{K} - \mathbf{K}_k \right\|_F^2$

# The Prototype Model

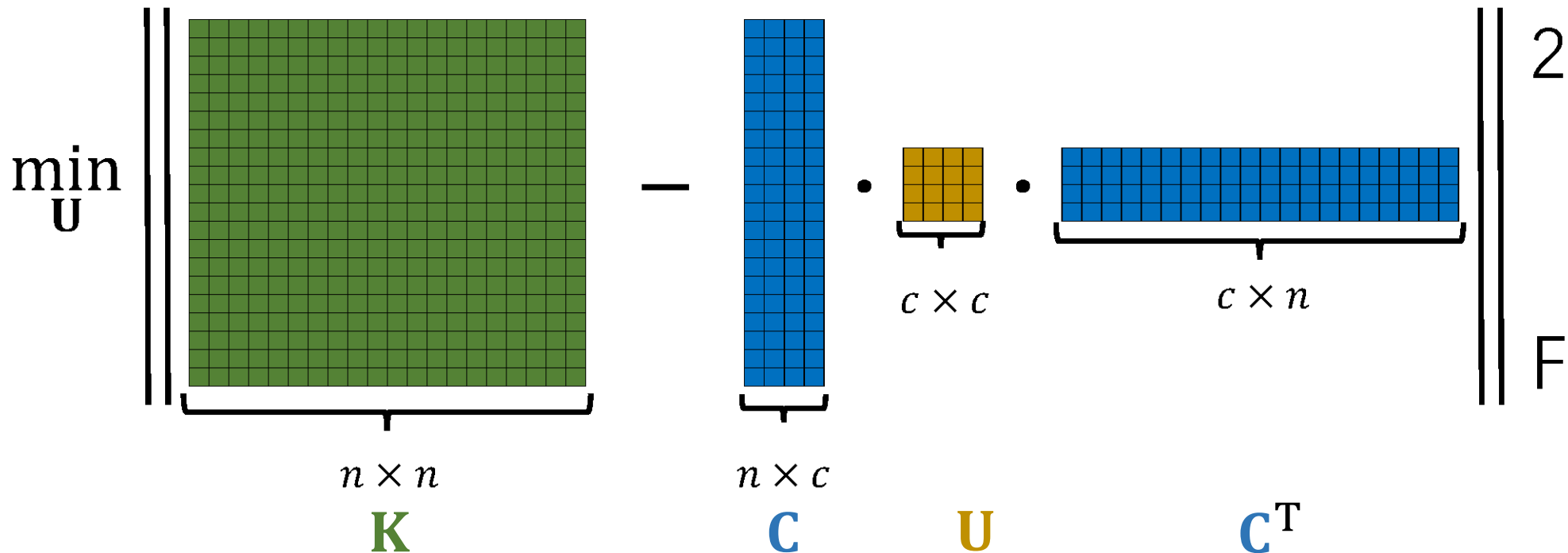
- Limitations

- $\mathbf{U}^* = \mathbf{C}^+ \mathbf{K} (\mathbf{C}^+)^T$
- Time cost is  $O(n^2 c)$
- Requires observing the whole of  $\mathbf{K}$

# The Prototype Model

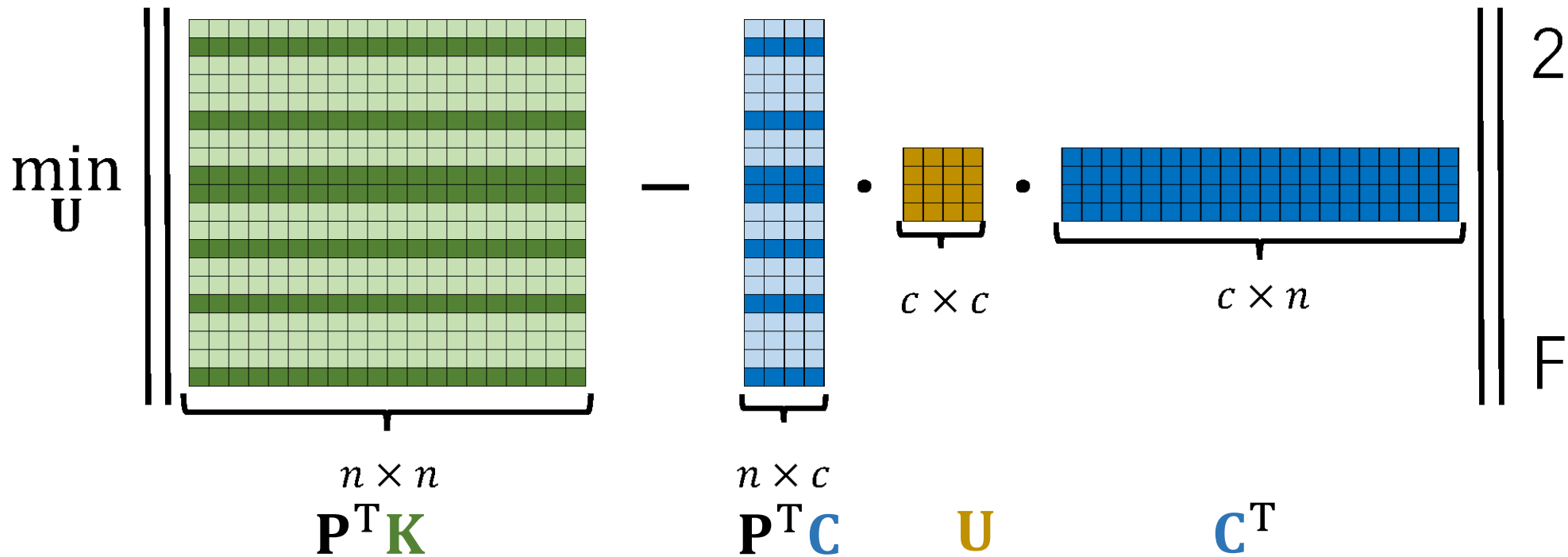
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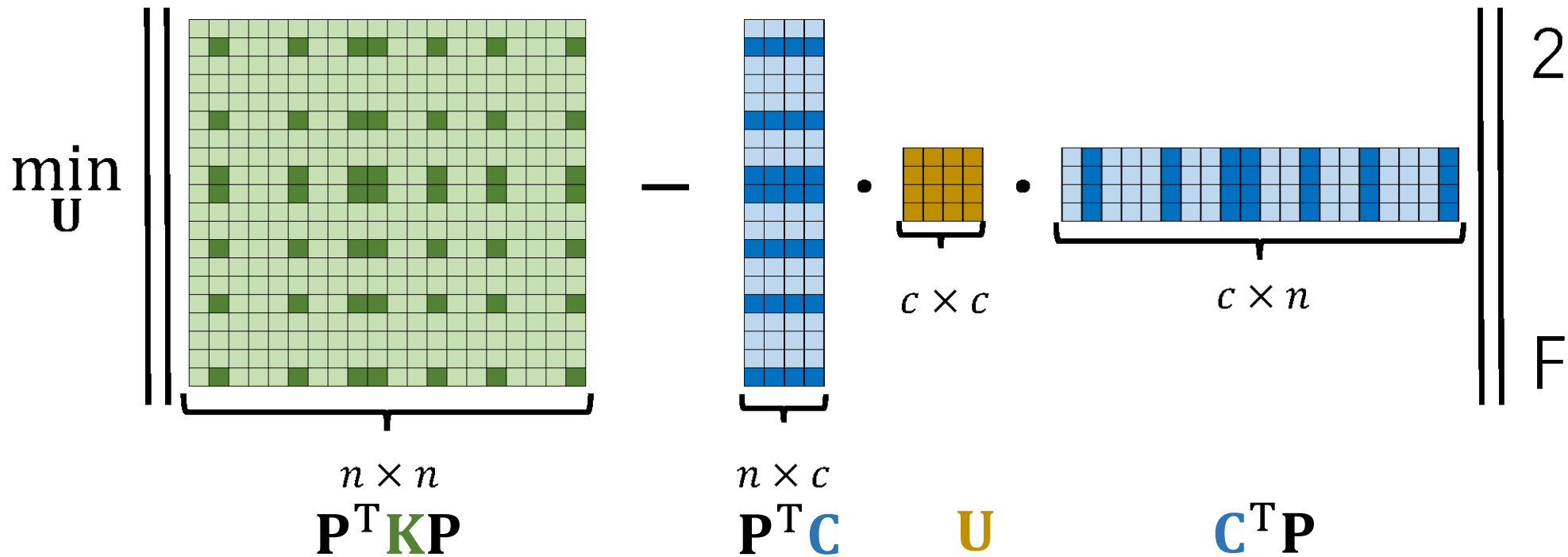
# The Fast Model

- Column/row selection
  - Form  $\mathbf{P}^T \mathbf{K} \mathbf{P}$  and  $\mathbf{P}^T \mathbf{C}$



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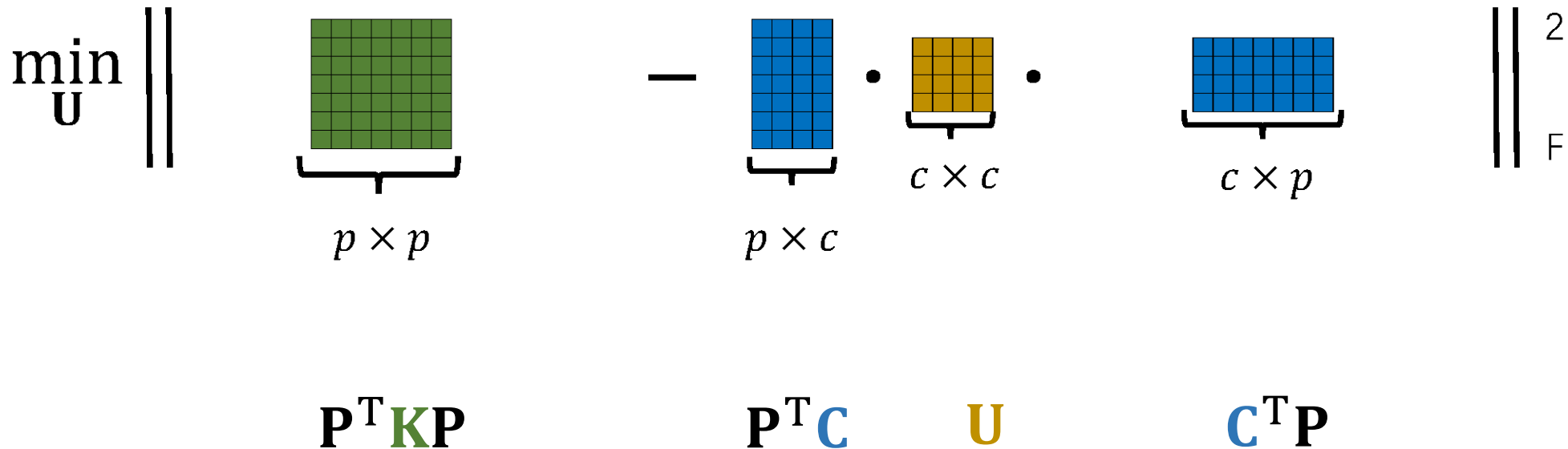




# The Fast Model

- $\mathbf{K} \approx \mathbf{C} \tilde{\mathbf{U}} \mathbf{C}^T$ , where

$$\tilde{\mathbf{U}} = \underset{\mathbf{U}}{\operatorname{argmin}} \left\| \mathbf{P}^T (\mathbf{K} - \mathbf{C} \mathbf{U} \mathbf{C}^T) \mathbf{P} \right\|_F^2.$$



# The Fast Model

- Prototype model:  $\mathbf{U}^* = \underset{\mathbf{U}}{\operatorname{argmin}} \left\| \mathbf{K} - \mathbf{C}\mathbf{U}\mathbf{C}^T \right\|_F^2 = \mathbf{C}^\dagger \mathbf{K} (\mathbf{C}^\dagger)^T$
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- Theory
  - $p = O\left(\sqrt{\frac{nc}{\epsilon}}\right)$
  - $\mathbf{P}$  is column selection matrix (according to the row leverage scores of  $\mathbf{C}$ )
  - Then  $\left\| \mathbf{K} - \mathbf{C}\tilde{\mathbf{U}}\mathbf{C}^T \right\|_F^2 \leq (1 + \epsilon) \left\| \mathbf{K} - \mathbf{C}\mathbf{U}^*\mathbf{C}^T \right\|_F^2$

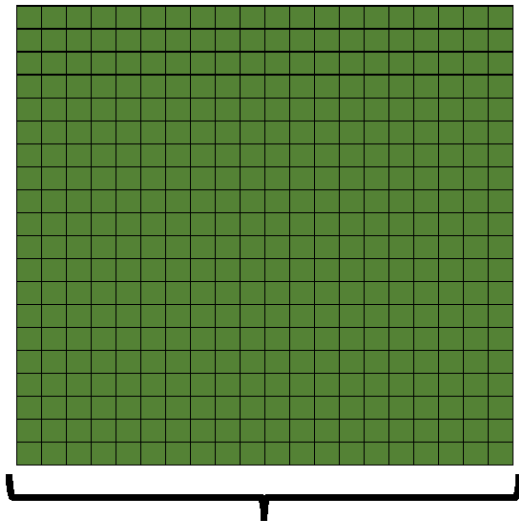
The faster model is nearly as good as the prototype model!

# The Fast Model

- Prototype model:  $\mathbf{U}^* = \underset{\mathbf{U}}{\operatorname{argmin}} \left\| \mathbf{K} - \mathbf{C}\mathbf{U}\mathbf{C}^T \right\|_F^2 = \mathbf{C}^\dagger \mathbf{K} (\mathbf{C}^\dagger)^T$
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- Overall time cost:  $O(p^2c + nc^2) = O(nc^3/\epsilon)$

linear in  $n$

# The Nystrom Method

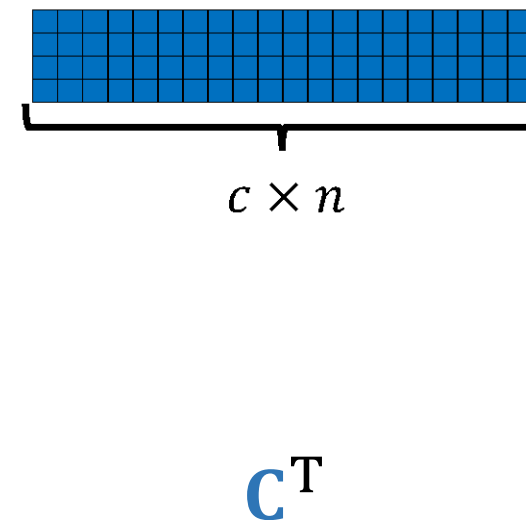
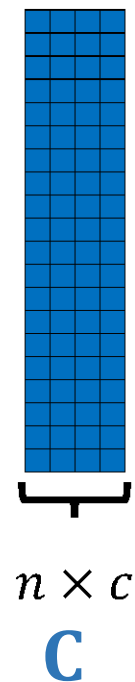
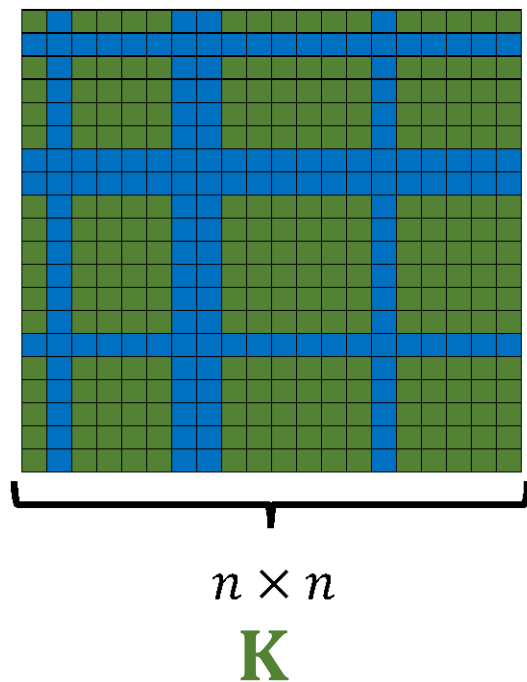


$n \times n$

**K**

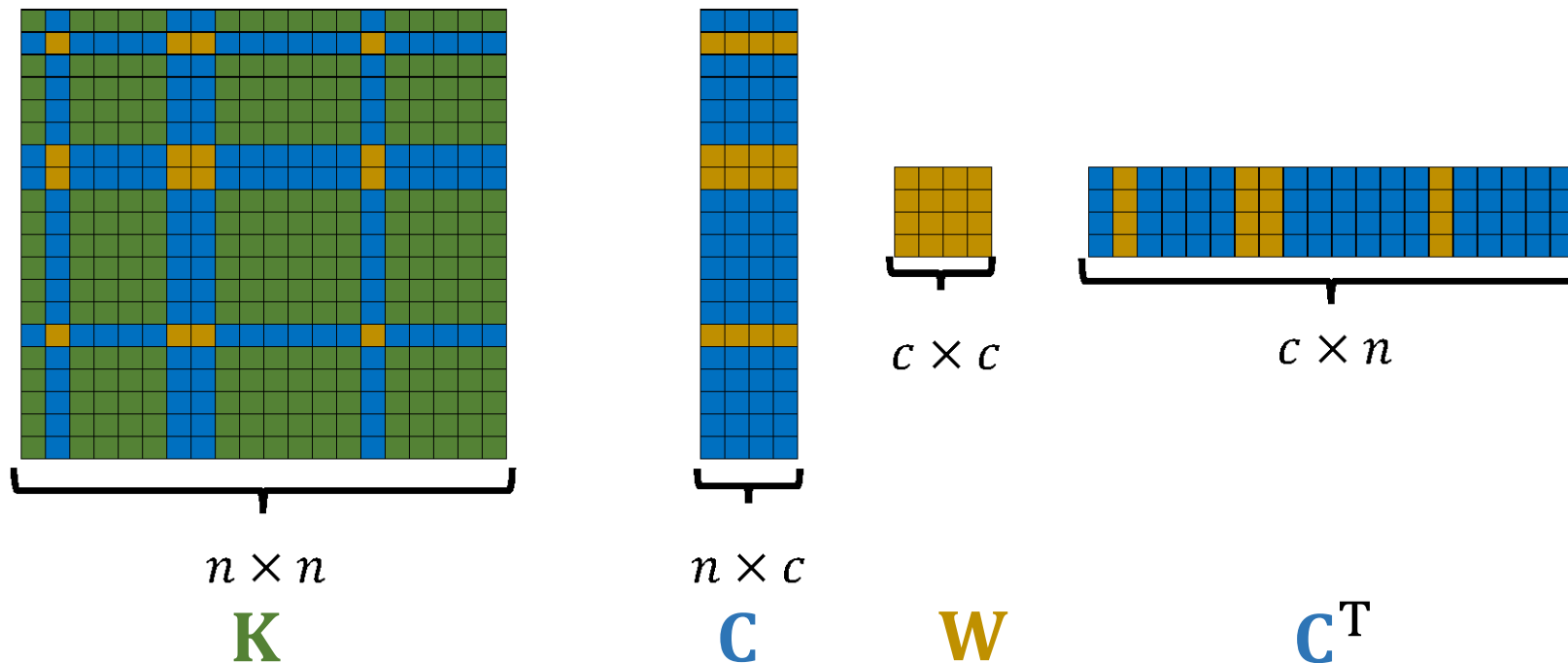
# The Nystrom Method

- $\mathbf{S}$  ( $n \times c$ ): column selection matrix
- $\mathbf{C} = \mathbf{KS}$  ( $n \times c$ )



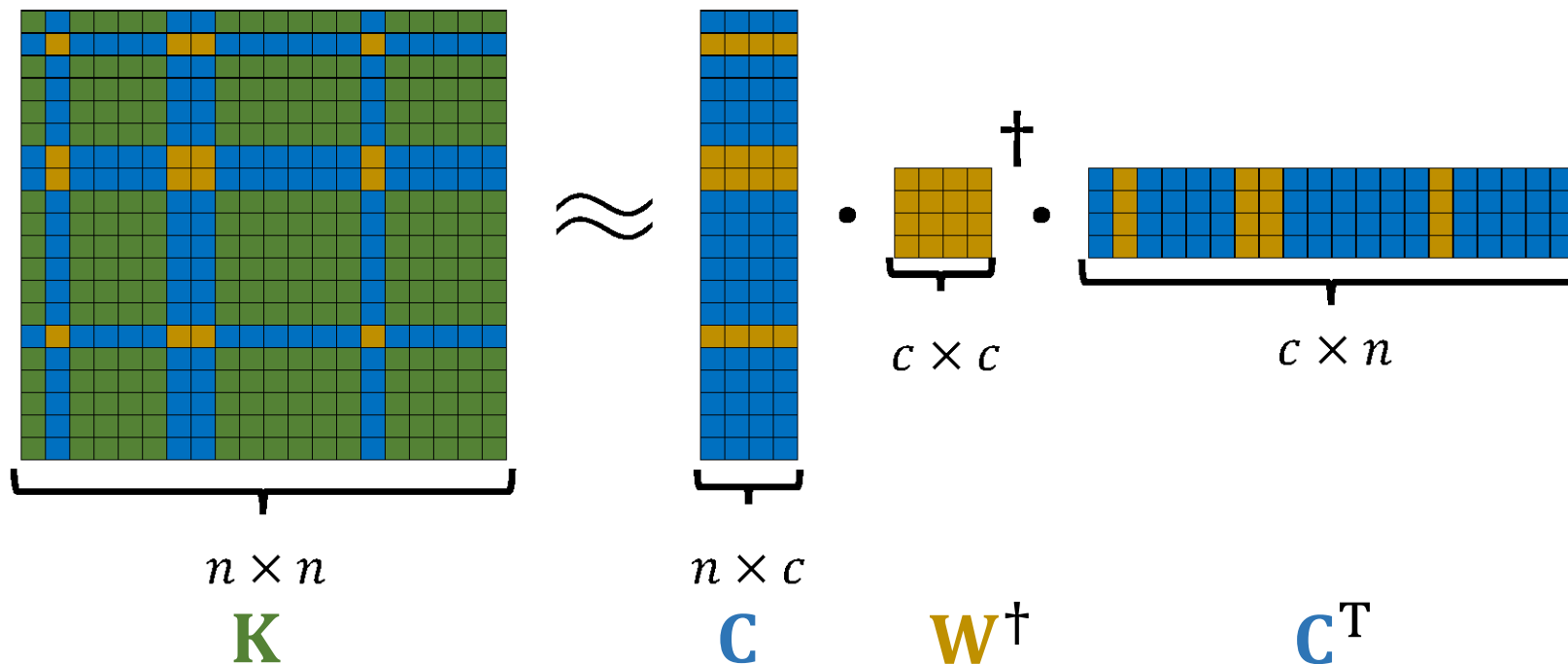
# The Nystrom Method

- $\mathbf{S}$  ( $n \times c$ ): column selection matrix
- $\mathbf{C} = \mathbf{KS}$  ( $n \times c$ ),  $\mathbf{W} = \mathbf{S}^T \mathbf{KS} = \mathbf{S}^T \mathbf{C}$  ( $c \times c$ )



# The Nyström Method

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- New explanation:
  - Recall the fast model:  $\tilde{\mathbf{X}} = \underset{\mathbf{X}}{\operatorname{argmin}} \left\| \mathbf{P}^T (\mathbf{K} - \mathbf{C} \mathbf{X} \mathbf{C}^T) \mathbf{P} \right\|_{\mathbf{F}}^2$

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- Setting  $\mathbf{P} = \mathbf{S}$ , then

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- The Nystrom method is special instance of the fast model.
- It is approximate solution to the prototype model

# The Nystrom Method

- Cost
  - Time:  $O(nc^2)$
  - Memory:  $O(nc)$



# The Nystrom Method

- Cost

- Time:  $O(nc^2)$

- Memory:  $O(nc)$

Very efficient!

# The Nystrom Method

- Cost
  - Time:  $O(nc^2)$
  - Memory:  $O(nc)$
- Error bound: weak

Very efficient!

# Comparisons

- $\mathbf{C} = \mathbf{KS} \in \mathbb{R}^{n \times c}$ ,  $\mathbf{W} = \mathbf{S}^T \mathbf{KS} = \mathbf{S}^T \mathbf{C} \in \mathbb{R}^{c \times c}$
- SPSD matrix approximation:  $\mathbf{K} \approx \mathbf{CUC}^T$ 
  - The prototype model:  $\mathbf{U} = \mathbf{C}^\dagger \mathbf{K} (\mathbf{C}^\dagger)^T$
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  - The Nyström method:  $\mathbf{U} = \mathbf{W}^\dagger$

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  - The Nyström method:  $\mathbf{U} = \mathbf{W}^\dagger$

When  $\mathbf{P} = \mathbf{I}_n$ , the prototype model  $\Leftrightarrow$  the fast model

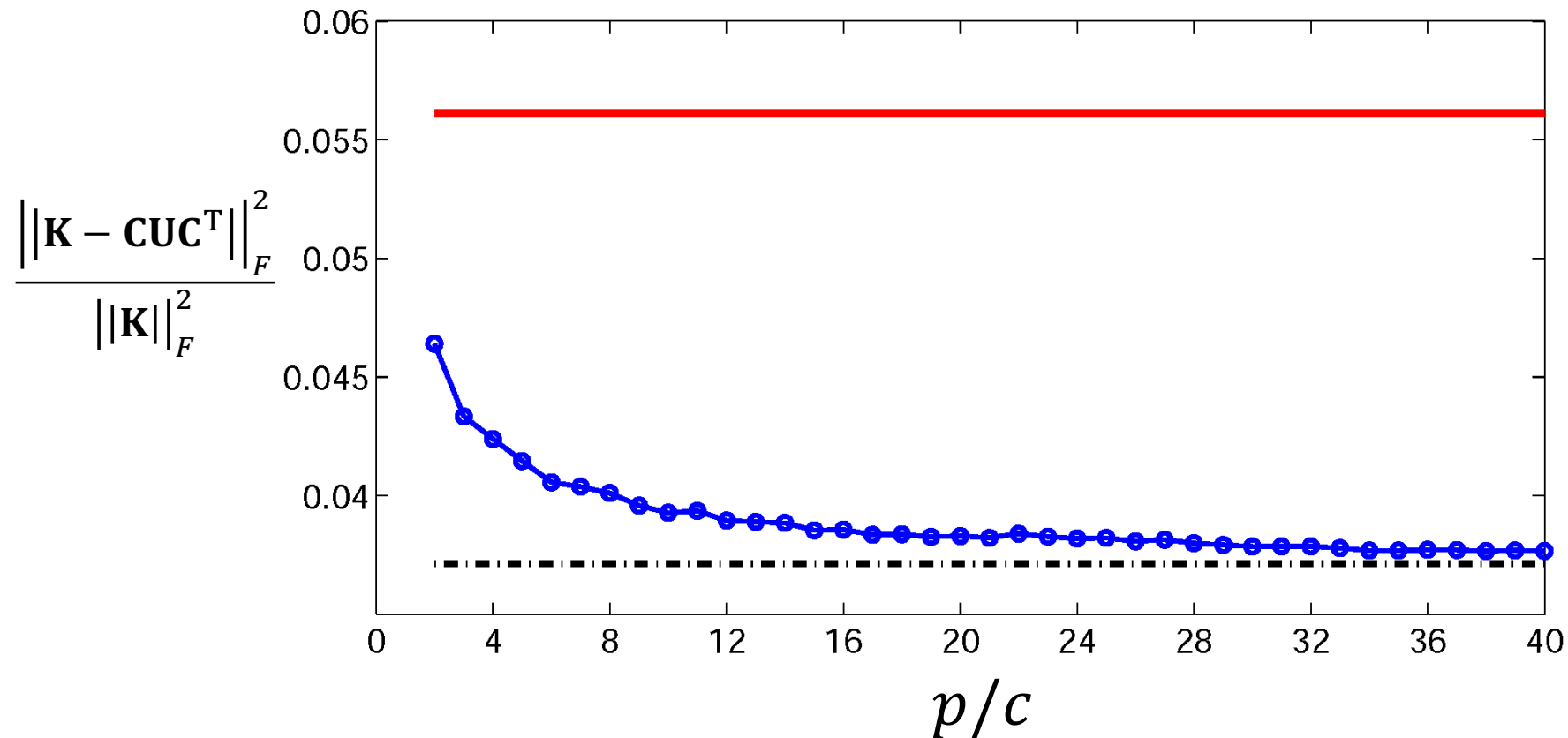
# Comparisons

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- SPSD matrix approximation:  $\mathbf{K} \approx \mathbf{C}\mathbf{U}\mathbf{C}^T$ 
  - The prototype model:  $\mathbf{U} = \mathbf{C}^\dagger \mathbf{K} (\mathbf{C}^\dagger)^T$
  - The fast model:  $\mathbf{U} = (\mathbf{P}^T \mathbf{C})^\dagger (\mathbf{P}^T \mathbf{K} \mathbf{P}) (\mathbf{C}^T \mathbf{P})^\dagger$
  - The Nyström method:  $\mathbf{U} = \mathbf{W}^\dagger$

When  $\mathbf{P} = \mathbf{S}$ , the Nyström method  $\Leftrightarrow$  the fast model

# Comparisons

- $c = 150, n = 100c$ , vary  $p$  from  $2c$  to  $40c$



**The Nyström Method**  
 $O(nc^2)$  time

**The Fast Model**  
 $O(nc^2 + p^2c)$  time

**The Prototype Model**  
 $O(n^2c)$  time

# Conclusions

- Motivations
  - Avoid forming the kernel matrix
  - Avoid inversion/decomposition
- Prototype model, fast model, Nystrom
  - They have connections
  - The fast model and Nystrom are practical